

The bandwidth conjecture for 3-colourable graphs

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(with Mathias Schacht & Anusch Taraz)

Outline

Let G be a graph.

We are interested in sufficient degree conditions on G for the appearance of a given graph H in G .

- 1 Introduction / Preliminary Results
- 2 The Bandwidth Conjecture
- 3 The Main Result
- 4 Outline of the Proof
- 5 Summary

Small Graphs

Theorem (follows from Turán 1941)

If $\delta(G) > \frac{r-2}{r-1}n$ then $K_r \subseteq G$.

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Theorem (follows from Erdős, Stone; Simonovits 1946; 1966)

For all graphs H and constants γ there is an n_0 s.t. for all graphs $|G| = n \geq n_0$ the following holds.

If $\delta(G) \geq \left(\frac{r-2}{r-1} + \gamma\right)n$ then $H \subseteq G$, where $r = \chi(H)$.

From Small Graphs to Big Graphs

- A graph H of constant size is forced in G , when

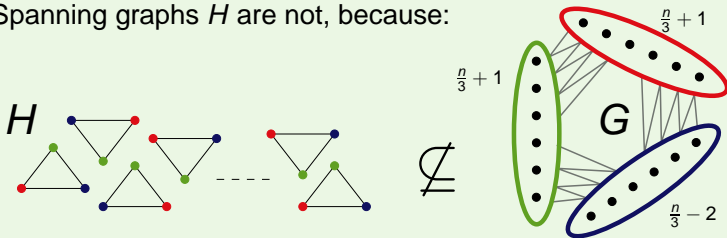
$$\delta(G) \geq \left(\frac{\chi(H) - 2}{\chi(H) - 1} + \gamma \right) n.$$

From Small Graphs to Big Graphs

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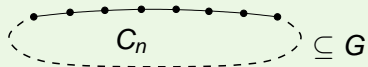
$$\delta(G) \geq \left(\frac{\chi(H) - 2}{\chi(H) - 1} + \gamma \right) n.$$

- Spanning graphs H are not, because:



Big Graphs

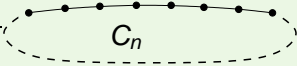
$$\blacksquare \delta(G) \geq \frac{1}{2}n$$

 \Rightarrow 

DIRAC'52

Big Graphs

■ $\delta(G) \geq \frac{1}{2}n \Rightarrow C_n \subseteq G$



The diagram shows a cycle graph C_n consisting of 8 vertices connected in a cycle. The vertices are arranged in a slightly curved line. A dashed oval encloses the entire cycle, and the label C_n is placed inside the oval. To the right of the oval is the subset symbol \subseteq followed by G .

DIRAC'52

■ $\delta(G) \geq \frac{r-1}{r}n \Rightarrow \frac{n}{r}$ disj. copies of $K_r \subseteq G$.

HAJNAL, SZEMERÉDI'69

Big Graphs

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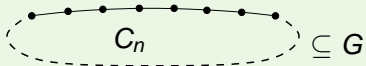
DIRAC'52

■ $\delta(G) \geq \frac{2}{3}n \Rightarrow K_3^* \subseteq G$

HAJNAL, SZEMERÉDI'69

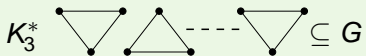
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DIRAC'52

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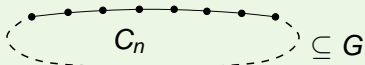
■ $\delta(G) \geq \frac{r-1}{r}n \Rightarrow (C_n)^r \subseteq G.$

FAN, KIERSTEAD '95

KOMLÓS, SÁRKÖZY, AND SZEMERÉDI '98

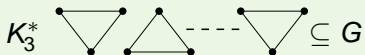
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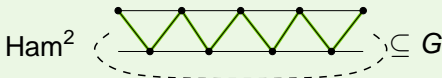
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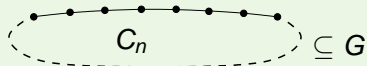


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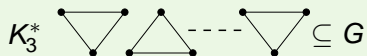
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KOMLÓS, SÁRKÖZY, AND SZEMERÉDI '98

■ other results: trees, F -factors, planar graphs . . .

KOMLÓS, SÁRKÖZY, AND SZEMERÉDI '95

ALON, YUSTER '96

KÜHN, OSTHUS, TARAZ '05

A Generalization

Conjecture (naïve)

*For all Δ and γ there is an n_0 s.t. for all $|H| = |G| = n \geq n_0$:
If $\Delta(H) \leq \Delta$ and $\delta(G) \geq \left(\frac{r-1}{r} + \gamma\right) n$ then $H \subseteq G$
where $r = \chi(H)$.*

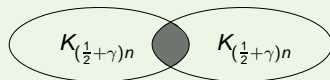
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Counterexample:

- H : random bipartite graph with $\Delta(H) \leq \Delta$.
- G : two cliques of size $\left(\frac{1}{2} + \gamma\right) n$ sharing $2\gamma n$ vertices.



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Bandwidth:

- $\text{bw}(G) \leq b$ if there is a labelling of $V(G)$ by $1, \dots, n$ s.t. for all $\{i, j\} \in E(G)$ we have $|i - j| \leq b$.



A Generalization

Conjecture (Bollobás & Komlós)

For all Δ and γ there are β and n_0 s.t. for all $|H| = |G| = n \geq n_0$:
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The 3-chromatic Case

Theorem

*For all Δ and γ there are β and n_0 s.t. for all $|H| = |G| = n \geq n_0$:
If $\Delta(H) \leq \Delta$, $\text{bw}(H) \leq \beta n$, $\chi(H) = 3$, and $\delta(G) \geq (\frac{2}{3} + \gamma) n$ then
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- The 2-chromatic case was announced by Abbasi (1998).

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- The proof is constructive.
- For the proof we use the regularity and the blow-up lemma.

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- 3 Take care of edges between the different parts of H with the help of (a variant of) the **embedding lemma**.

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- 2 Prepare H by exploiting its structure and find an assignment of $V(H)$ to the partition classes of G . (Lemma for H)
- 3 Take care of edges between the different parts of H with the help of (a variant of) the embedding lemma.
- 4 Complete the embedding with the blow-up lemma.

The Lemma for G

Lemma (Lemma for G)

Let G be a suff. large graph with $\delta(G) \geq (\frac{2}{3} + \gamma)n$. There is a **regular partition** $V(G) = V'_1 \dot{\cup} \dots \dot{\cup} V'_k$ of G with **reduced graph R** with the following properties:

- $\delta(R) \geq (\frac{2}{3} + \gamma/2)|V(R)|$ and thus $R \supseteq \text{Ham}^2 \supseteq K_3^*$,
- $V'_1 \dot{\cup} \dots \dot{\cup} V'_k$ is super-regular on K_3^* ,
- The parts V'_i are almost of equal size.

Moreover, if we are asked to change the sizes of the parts V'_i by ξn , we can do so. (resulting partition: $V_1 \dot{\cup} \dots \dot{\cup} V_k$)

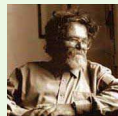
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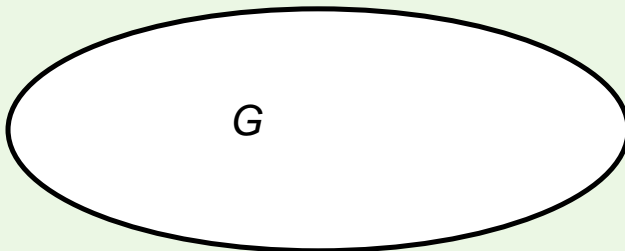
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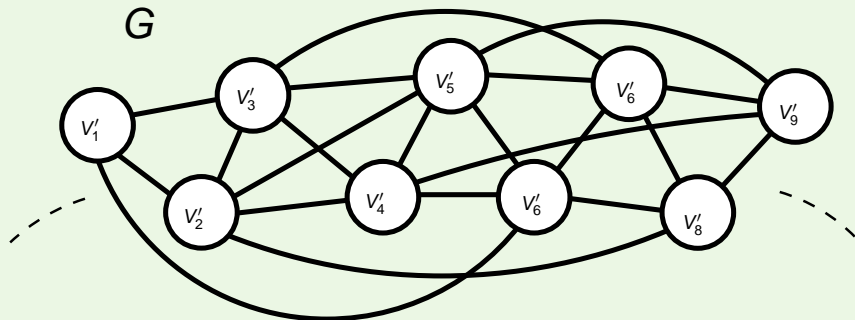
Moreover, if we are asked to change the sizes of the parts V'_i by ξn , we can do so. (resulting partition: $V_1 \dot{\cup} \dots \dot{\cup} V_k$)



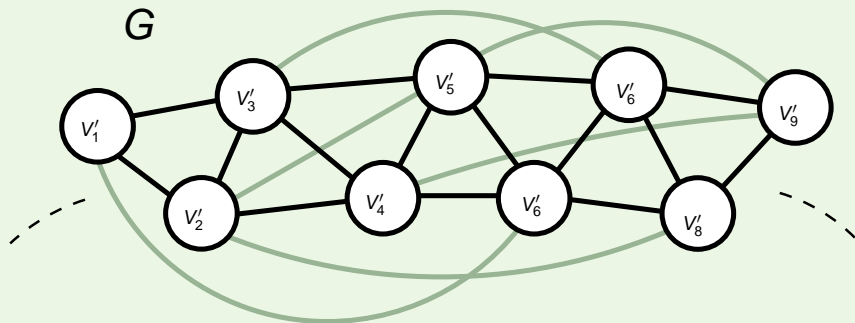
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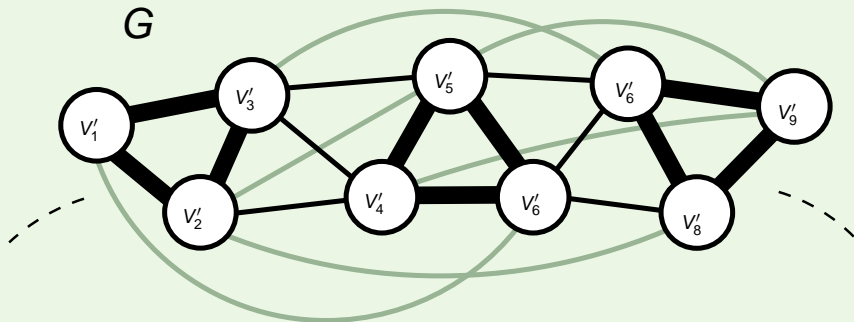


A regular partition of G with parts of almost equal size.

The Lemma for G 

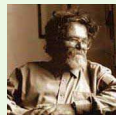
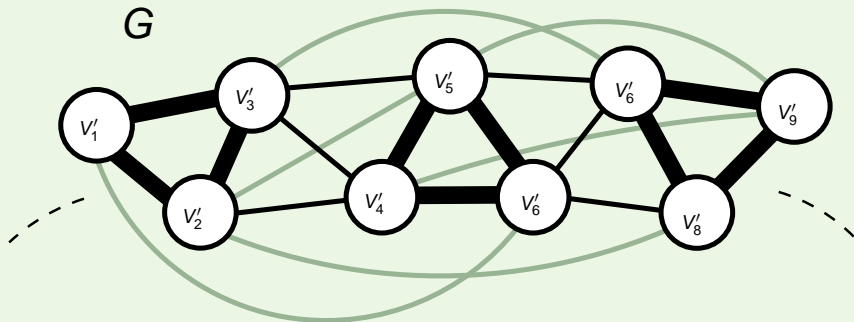
Since $\delta(R) \geq (\frac{2}{3} + \gamma/2)|V(R)|$ we have $R \supseteq \text{Ham}^2 \supseteq K_3^*$.

The Lemma for G

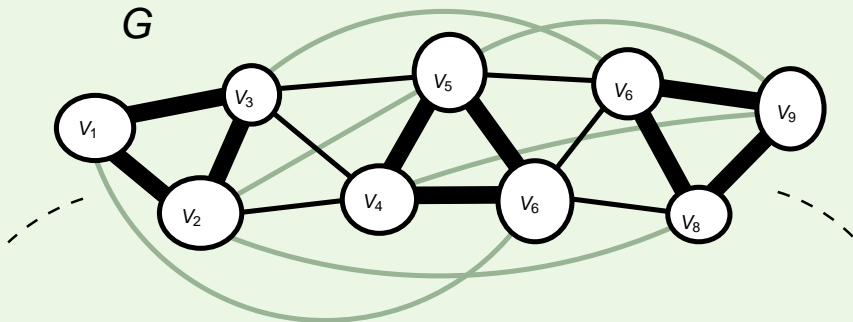


We make the partition super-regular on K_3^* .

The Lemma for G



The Lemma for G



We can change the partition by moving some vertices to obtain $V_1 \dot{\cup} \dots \dot{\cup} V_k$.

The Lemma for H

Lemma (Lemma for H)

Let H be a graph with $\Delta(H) \leq \Delta$, bandwidth $\leq \beta n$ and $\chi(H) \leq 3$, and let $R \supseteq \text{Ham}^2 \supseteq K_3^*$ be a graph with $\delta(R_k) > 2k/3$. Then there is a **homomorphism f from H to R** such that

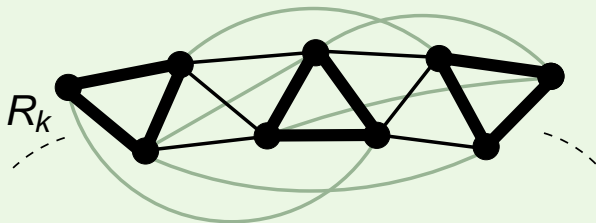
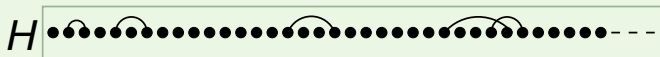
- f sends roughly the same number of vertices to each $v \in V(R)$,
- Almost all edges of H are mapped to edges of K_3^* .

Moreover, the set of **special vertices**

$$X := \{v \in V(H) : \text{some edge } uv \text{ is not mapped to } K_3^*\}$$

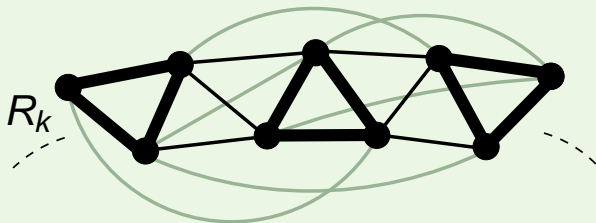
is very small.

The Lemma for H



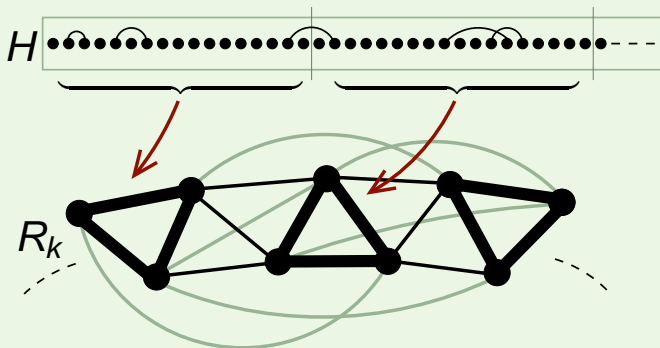
H is given in an order respecting the bandwidth bound.

The Lemma for H

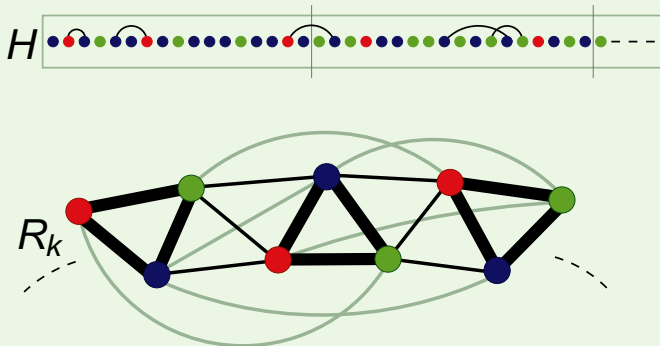


Idea: Cut H into pieces

The Lemma for H



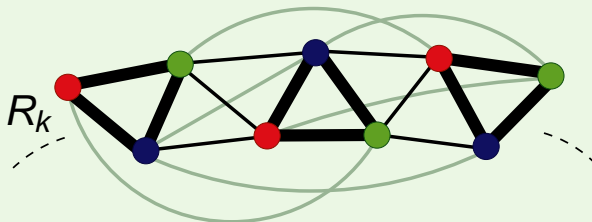
Idea: Cut H into pieces and map each piece to a triangle of K_3^* .

The Lemma for H 

This is possible, since $\chi(H) = 3$.

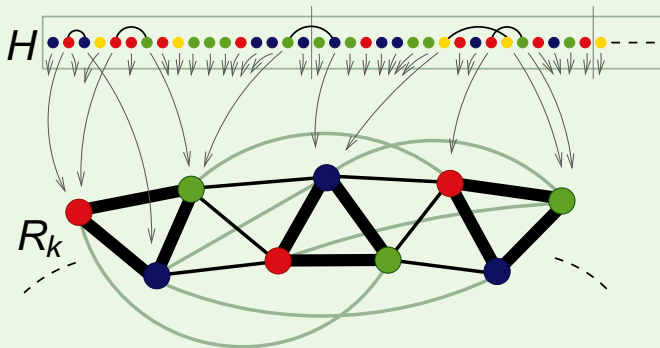
But the colour classes of H may vary in size a lot.

The Lemma for H



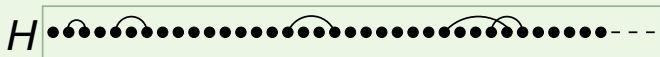
We can find a large subgraph of H with a balanced 3-colouring.

The Lemma for H



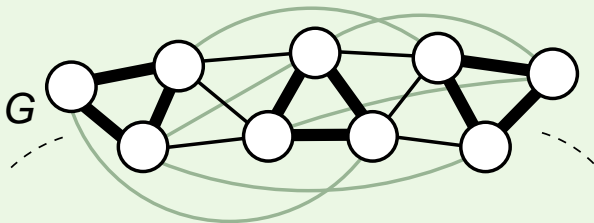
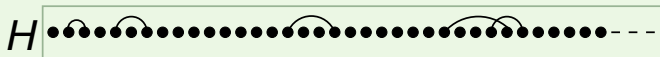
We map all edges of H to R , and most of them to K_3^* .

Proof of the Main Theorem



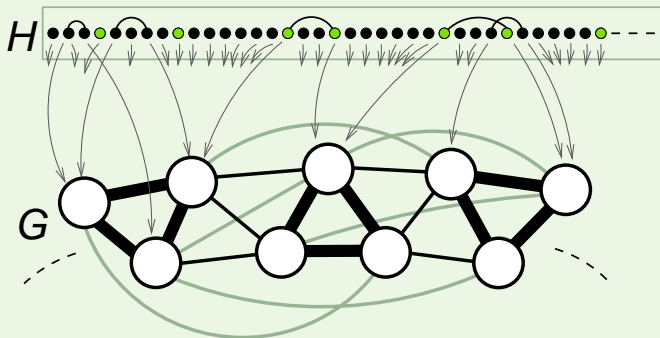
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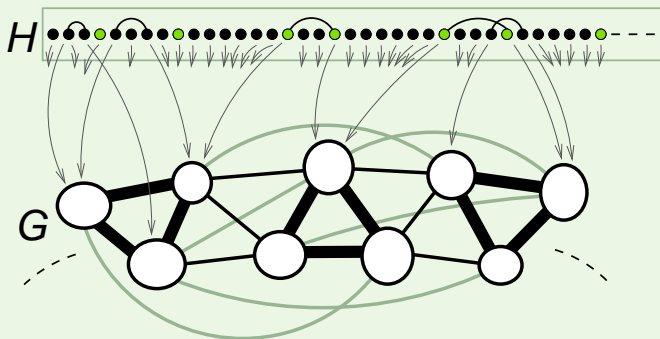
The **Lemma for G** constructs a regular partition $V_1' \dot{\cup} \dots \dot{\cup} V_k'$ of G .

Proof of the Main Theorem



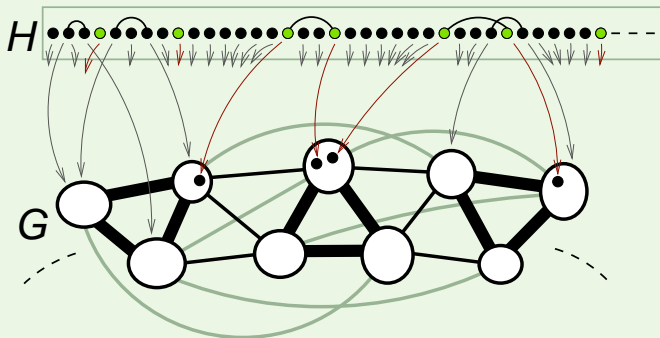
The **Lemma for H** constructs a homomorphism f and a set of special vertices X .

Proof of the Main Theorem



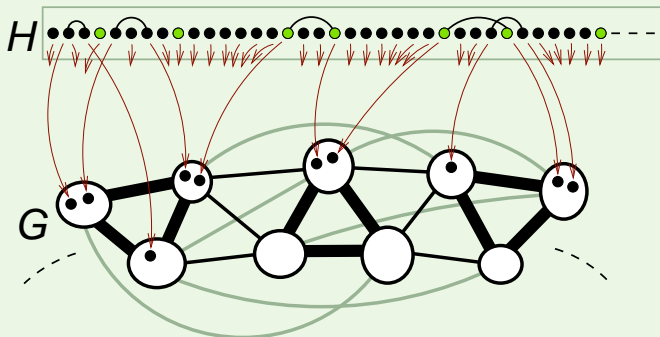
The **Lemma for G** adjusts the partition of G s.t. $|V_i| = f^{-1}(i)$.

Proof of the Main Theorem



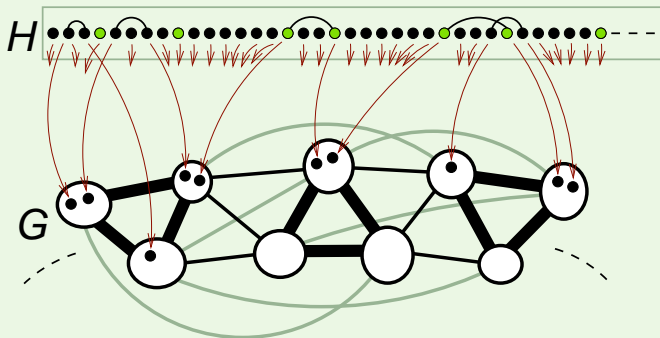
We embed the special vertices X into G using the **embedding lemma**.

Proof of the Main Theorem



We embed all other vertices using the **blow-up lemma**.

Proof of the Main Theorem



$$H \subseteq G$$

Summary/Outlook

- The bandwidth conjecture generalizes many known results about sufficient degree conditions for the appearance of spanning subgraphs.
- We proved the bandwidth conjecture for $\chi(H) = 3$.
- The proof does not (quite) carry over to $\chi(H) > 3$.
- What happens if we replace bandwidth by pathwidth / treewidth?