

# 47853 Packing and Covering: Lecture 1

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## 1 What is packing and covering?

### 1.1 Menger's theorem and its dual

Let  $G = (V, E)$  be a graph, and take distinct vertices  $s, t \in V$ .<sup>1</sup> An *st-path* is a minimal edge subset connecting  $s$  and  $t$ . An *st-cut* is an edge subset of the form

$$\delta(U) := \{e \in E : |e \cap U| = 1\}$$

where  $U \subseteq V$  satisfies  $U \cap \{s, t\} = \{s\}$ . We will refer to  $U$  and  $V - U$  as the *shores* of  $G$ . Notice that every *st-path* intersects every *st-cut*.

What is the maximum number of (pairwise) disjoint *st-paths*? In other words, how many *st-paths* can we *pack*?

**Theorem 1.1** (Menger 1927 [10]). *Let  $G = (V, E)$  be a graph, and take distinct vertices  $s, t \in V$ . Then the maximum number of disjoint *st-paths* is equal to the minimum cardinality of an *st-cut*.*

*Proof.* Every *st-path* intersects an *st-cut*, so the maximum number of disjoint *st-paths* is at most the minimum cardinality of an *st-cut*. We prove the other inequality by induction on  $|V| + |E| \geq 3$ . The result is obvious for  $|V| + |E| = 3$ . For the induction step, assume that  $|V| + |E| \geq 4$ . Let  $\tau$  be the minimum cardinality of an *st-cut*. We may assume that  $\tau \geq 1$ . We will find  $\tau$  disjoint *st-paths*.

**Claim 1.** *If an edge  $e$  does not appear in a minimum *st-cut*, then  $G$  has  $\tau$  disjoint *st-paths*.*

*Proof of Claim.* Notice that the cardinality of a minimum *st-cut* in  $G \setminus e$  is still  $\tau$ . As a result, the induction hypothesis implies the existence of  $\tau$  disjoint *st-paths* in  $G \setminus e$ , and therefore in  $G$ .  $\diamond$

We may therefore assume that every edge appears in a minimum *st-cut*. An *st-cut*  $\delta(U)$  is *trivial* if either  $|U| = 1$  or  $|V - U| = 1$ .

**Claim 2.** *If there is a minimum *st-cut* that is not trivial, then  $G$  has  $\tau$  disjoint *st-paths*.*

<sup>1</sup>We allow parallel edges but disallow loops, until further notice.

*Proof of Claim.* Let  $\delta(U), s \in U \subseteq V - \{t\}$  be a minimum  $st$ -cut that is non-trivial. Let  $G_1$  be the graph obtained from  $G$  by shrinking  $U$  to a single vertex  $s'$ , and let  $G_2$  be the graph obtained from  $G$  after shrinking  $V - U$  to a single vertex  $t'$ . Since  $\delta(U)$  is non-trivial, it follows that  $|V(G_i)| + |E(G_i)| < |V| + |E|$ , for each  $i \in [2]$ . We may therefore apply the induction hypothesis to  $G_1$  and  $G_2$ . Notice that  $\tau$  is still the minimum cardinality of an  $s't$ -cut in  $G_1$  and of an  $s't'$ -cut in  $G_2$ . Thus, by the induction hypothesis,  $G_1$  has  $\tau$  disjoint  $s't$ -paths and  $G_2$  has disjoint  $s't'$ -paths. Gluing these paths along the edges of  $\delta(U)$  gives us  $\tau$  disjoint  $st$ -paths in  $G$ .  $\diamond$

We may therefore assume that every minimum  $st$ -cut is trivial. Since every edge appears in a minimum  $st$ -cut, it follows that every edge has either  $s$  or  $t$  as an end. In this case,  $G$  has a special form and it is clear that  $\tau = \nu$  for this graph, thereby completing the induction step.  $\square$

On the other hand, how many  $st$ -cuts can we pack?

**Theorem 1.2.** *Let  $G = (V, E)$  be a connected graph  $G$ , and take distinct vertices  $s, t \in V$ . Then the maximum number of disjoint  $st$ -cuts is equal to the minimum cardinality of an  $st$ -path.*

*Proof.* Clearly, the maximum number of disjoint  $st$ -cuts is at most the minimum cardinality of an  $st$ -path. To prove the other inequality, let  $\tau \geq 1$  be the minimum cardinality of an  $st$ -path. We will find  $\tau$  disjoint  $st$ -cuts. Notice that  $\tau$  is equal to the distance between  $s$  and  $t$ . For each  $i \in \{0, 1, \dots, \tau - 1\}$ , let  $U_i$  be the set of vertices at distance at most  $i$  from  $s$ . Notice that  $\{s\} = U_0 \subsetneq U_1 \subsetneq \dots \subsetneq U_{\tau-1} \subseteq V - \{t\}$ . Our definition implies that  $\delta(U_0), \delta(U_1), \dots, \delta(U_{\tau-1})$  are disjoint  $st$ -cuts, as required.  $\square$

These results are two of many packing theorems. Just to mention a few, we will see some of these packing results:

- Lucchesi and Younger 1978 [9]: given a directed graph  $G$ , the maximum number of disjoint dicuts is equal to the minimum cardinality of a dijoin.
- Conjecture (Woodall 1978 [13]): given a directed graph  $G$ , the maximum number of disjoint dijoins is equal to the minimum cardinality of a dicut.
- Edmonds and Johnson 1973 [4]: given a graph  $G$  and even subset  $T$  of vertices, the maximum value of a fractional packing of  $T$ -joins is equal to the minimum cardinality of a  $T$ -cut.
- Guenin 2001 [7]: in a signed graph without an odd- $K_5$  minor, the maximum value of a fractional packing of odd circuits is equal to the minimum cardinality of a signature.

## 1.2 Dilworth's theorem and its dual

Take a partially ordered set  $(E, \leq)$ , that is, the following statements hold for all  $a, b, c \in E$ :

- $a \leq a$ ,

- if  $a \leq b$  and  $b \leq a$ , then  $a = b$ ,
- if  $a \leq b$  and  $b \leq c$ , then  $a \leq c$ .

We say that  $a, b$  are *comparable* if  $a \geq b$  or  $b \geq a$ ; otherwise they are *incomparable*. A *chain* is a set of pairwise comparable elements. An *antichain* is a set of pairwise incomparable elements. Notice that every antichain intersects every chain at most once.

What is the minimum number of (not necessarily disjoint) chains whose union is  $E$ ? That is, what is the least number of chains needed to *cover* the ground set?

**Theorem 1.3** (Dilworth 1950 [2]). *Let  $(E, \leq)$  be a partially ordered set. Then the minimum number of chains needed to cover  $E$  is equal to the maximum cardinality of an antichain.*

*Proof.* Since every chain intersects every antichain at most once, the minimum number of chains needed to cover  $E$  is greater than or equal to the maximum cardinality of an antichain. We will prove the other inequality by induction on  $|E|$ . The base case  $|E| = 1$  is obvious. For the induction step, assume that  $|E| \geq 2$ . Let  $\alpha$  be the maximum cardinality of an antichain. We will find  $\alpha$  chains covering  $E$ . If  $\alpha = |E|$ , then we are clearly done. Otherwise,  $\alpha < |E|$ , implying in turn that there is a chain  $\{a, b\}$  where  $a$  is a minimal element and  $b$  is a maximal element. Let  $E' := E - \{a, b\}$ .

**Claim.** *If the maximum cardinality of an antichain of  $(E', \leq)$  is  $\alpha - 1$ , then there are  $\alpha$  chains covering  $E$ .*

*Proof of Claim.* By the induction hypothesis, there are  $\alpha - 1$  chains of  $E'$  covering  $E - \{a, b\}$ . Together with  $\{a, b\}$ , we get a covering of  $E$  using  $\alpha$  chains.  $\diamond$

We may therefore assume that  $E'$  has an antichain  $A$  such that  $|A| = \alpha$ . Let

$$E^+ := A \cup \{x \in E - A : x \geq z \text{ for some } z \in A\}$$

$$E^- := A \cup \{y \in E - A : y \leq z \text{ for some } z \in A\}.$$

Since  $A$  is an antichain,  $E^+ \cap E^- = A$ , and since it is a maximum antichain,  $E^+ \cup E^- = E$ . As  $a$  is minimal and  $a \notin A$ , it follows that  $a \notin E^+$ . As  $b$  is maximal and  $b \notin A$ , we get that  $b \notin E^-$ . In particular,  $|E^+|, |E^-| < |E|$ . Thus, by the induction hypothesis,  $E^+$  has  $\alpha$  chains covering it, and  $E^-$  has  $\alpha$  chains covering it. Gluing these chains together, we get  $\alpha$  chains covering  $E^+ \cup E^- = E$ , thereby completing the induction step.  $\square$

On the other hand, what is the least number of antichains needed to cover the ground set?

**Theorem 1.4.** *Let  $(E, \leq)$  be a partially ordered set. Then the minimum number of antichains needed to cover  $E$  is equal to the maximum cardinality of a chain.*

*Proof.* Clearly, the minimum number of antichains needed to cover  $E$  is greater than or equal to the maximum cardinality of a chain. To prove the other inequality, let  $\alpha$  denote the maximum cardinality of a chain. We will find  $\alpha$  antichains whose union is  $E$ . Let  $A_1$  denote the set of all minimal elements of  $E$ . For each  $i \geq 2$ , let  $A_i$  denote the set of all minimal elements of  $E - (A_1 \cup \dots \cup A_{i-1})$ . Observe that

- $E = \bigcup_{i \geq 1} A_i$ ,
- each  $A_i$  is an antichain,
- if  $i \geq 2$  and  $a \in A_i$ , then there is a  $b \in A_{i-1}$  such that  $a \geq b$ , and so
- if  $A_i \neq \emptyset$ , then there is a chain of cardinality  $i$ .

As a result, since  $\alpha$  is the maximum cardinality of a chain, it follows that  $\emptyset = A_{\alpha+1} = A_{\alpha+2} = \dots$ . Thus,  $E$  is the union of the  $\alpha$  antichains  $A_1, \dots, A_\alpha$ , as required.  $\square$

These results are two of many covering results. To name a few:

- Kőnig 1931 [8]: In a bipartite graph, the minimum number of colors needed for a proper edge-coloring is equal to the maximum degree of a vertex.
- Gallai 1962 [6], Surányi 1968 [12]: In a chordal graph, the minimum number of cliques needed to cover the vertices is equal to the maximum cardinality of a stable set.
- Sachs 1970 [11]: In a chordal graph, the minimum number of colors needed for a proper vertex-coloring is equal to the maximum cardinality of a clique.
- Chudnovsky, Robertson, Seymour and Thomas 2006 [1]: In a graph without an odd hole or an odd hole complement, the minimum number of cliques needed to cover the vertices is equal to the maximum cardinality of a stable set.

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